

Forwarding with in-packet Bloom Filters

T-110.6120 9.10.2012

Jimmy Kjällman Ericsson Research, NomadicLab



Background



General Starting Points

- New Future Internet architecture
- Focus on long-term research
 - With feedback to short-term work
- Clean slate approach
 - Reconsidering old assumptions
- Redesigning the Internet architecture
 - Considering both technical and socio-economic aspects

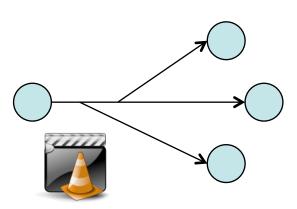


- Information-Centric Networking
 - Various projects around the world



Choices and Goals (and Constraints)

- Information-centric
 - Not host centric
- Publish/subscribe
 - Instead of send/receive
- Identify information
 - No (global) node addresses
- Secure and efficient networking
 - DDoS protection, multicast, ...





Projects

EU FP7 PSIRP 2008-2010
 Publish/Subscribe Internet Routing Paradigm



EU FP7 PURSUIT 2010-2013
 Publish/Subscribe Internet Technology



• ICT SHOK FI WP3 2008-2012

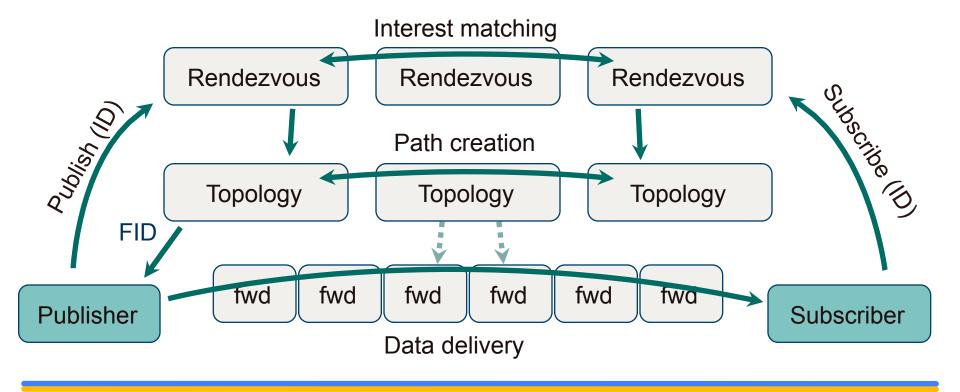




PSIRP/PURSUIT Basic Architectural functions



- Rendezvous matching publish and subscribe events
- Topology network topology knowledge, path computation
- Forwarding fast data delivery





Ideas about Forwarding

- Need for a new forwarding mechanism in PSIRP
- Some requirements
 - Multicast support
 - Security (receiver in control, DDoS protection)
 - Efficiency
- One of the initial ideas: MPLS-like labels
- Another idea: Bloom filters
 - Very little state and signaling required, native multicast support, no global addressing, path not revealed, no routing tables and lookups, no pushing/popping, ...



<u>Li</u>ne Speed <u>P</u>ublish/<u>S</u>ubscribe <u>I</u>nter-<u>N</u>etworking

 Petri Jokela^(*), András Zahemszky, Christian Esteve, Somaya Arianfar, and Pekka Nikander, "LIPSIN: Line speed Publish/Subscribe Inter-Networking", ACM SIGCOMM 2009



Bloom filters – Burton Howard Bloom, 1970



Bloom filters

- Probabilistic data structure, space efficient
- Used to test if an element has been added to a set

10-bit Bloom Filter

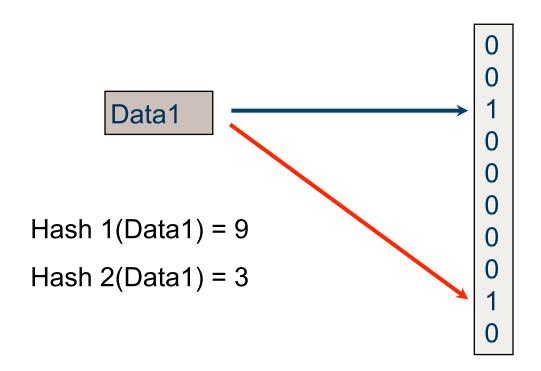
Hash 1 Hash 2



Bloom filters: Inserting items

 Hash the data k times, get index values, and set the bits





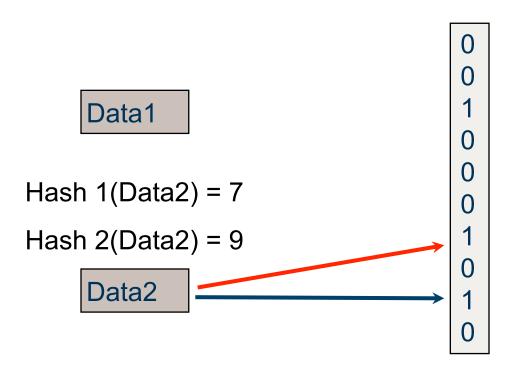
Hash 1 Hash 2



Bloom filters: Inserting items

 Hash the data k times, get index values, and set the bits





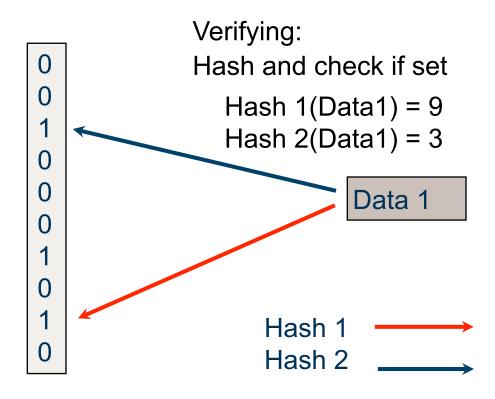
Hash 1 Hash 2



Bloom filters: Verifying (positive)

All corresponding bits have been set → positive response

10-bit Bloom Filter

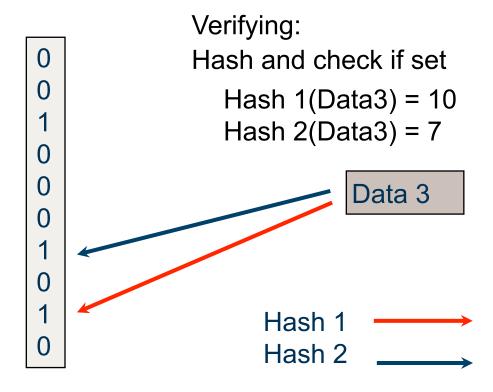




Bloom filters: Verifying (negative)

Some bits do not match → negative response



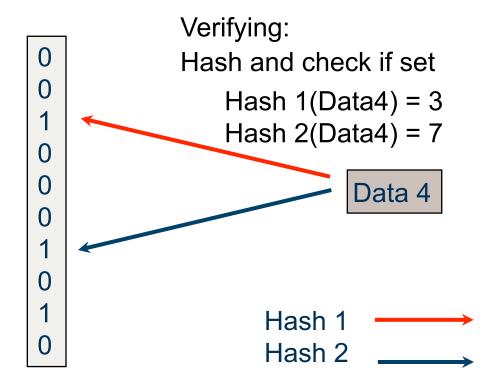




Bloom filters: False positives

Bits match the BF although "Data 4" was never added

10-bit Bloom Filter



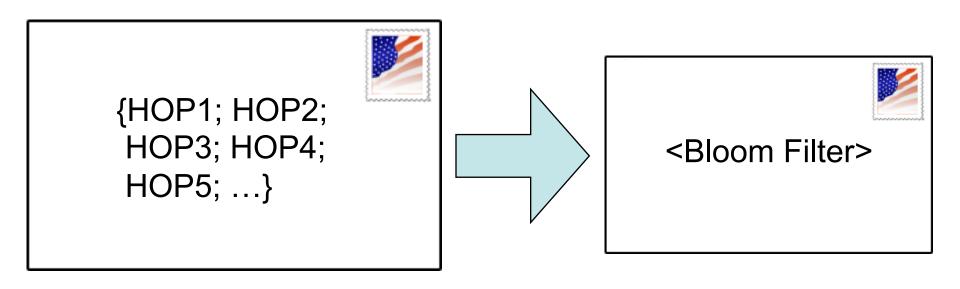


In-packet bloom filters- zFilters



Forwarding with zFilters

- Source routing
- Explicitly enumerating all hops requires a lot of space
 - so instead we encode this information into a Bloom filter

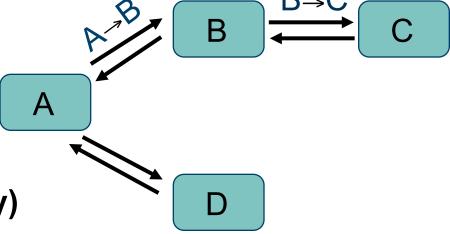


No names for nodes

 Each link is identified with a unidirectional (outgoing) Link ID

Link IDs

- No hashing required,
 generate the 1-bits
 otherwise (e.g. randomly)
- Size e.g. 256 bits of which5 bits set to 1
 - 2 x the size of an IPv6 addr
 - Statistically unique



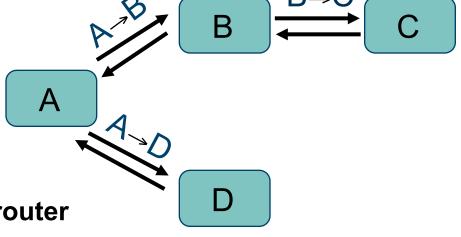
Link IDs and zFilters

A->B

B->C

A->C

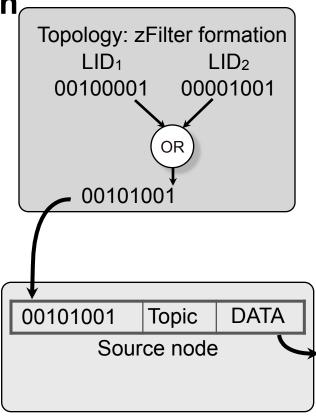
- Strict source routing
 - Create a path, collect all Link IDs
 - Include (OR) all path's/tree's
 Link IDs into a Bloom filter
- Multicast support
 - Include multiple outgoing links from one router
- Stateless (almost)
 - Only Link IDs stored on the router
- Packet forwarding
 - Always to the correct destination
 - False positives possible





Topology manager's role

- Needs (intra-)network link information
 - Topology and Link IDs
 - E.g., OSPF, PCE
- Computes paths on request
 - Creates the zFilter using the Link ID information
 - Gives the zFilter to the source node
 - (Source adds zFilter to outgoing data packets)





Forwarding decision

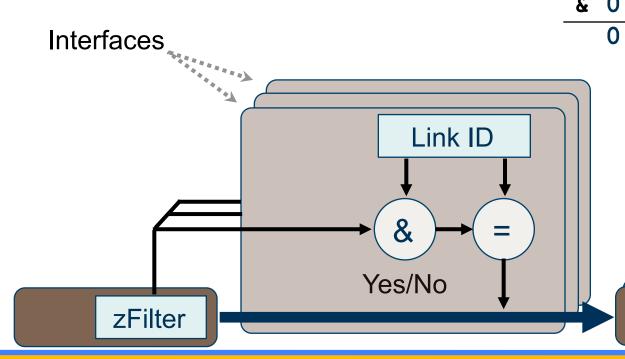
- Forwarding decision based on binary AND and a comparison
 - zFilter in the packet matched with all outgoing Link IDs
 - Forward if: zFilter AND LID = LID

```
(\Leftrightarrow (zFilter\ AND\ LID)\ XOR\ LID = 0)
```

 1
 1
 0
 0
 0
 1
 1
 0
 1

 &
 0
 1
 0
 0
 1
 0
 0
 1

 0
 1
 0
 0
 0
 1
 0
 0
 1

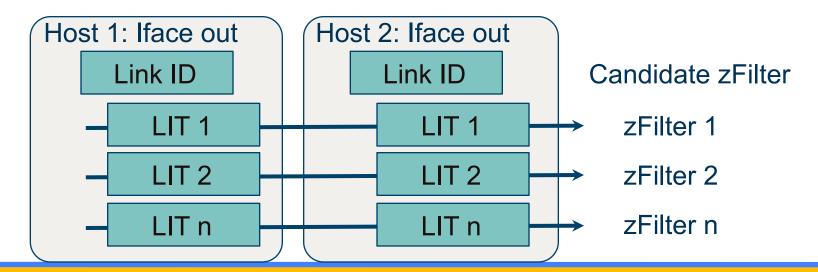


zFilter



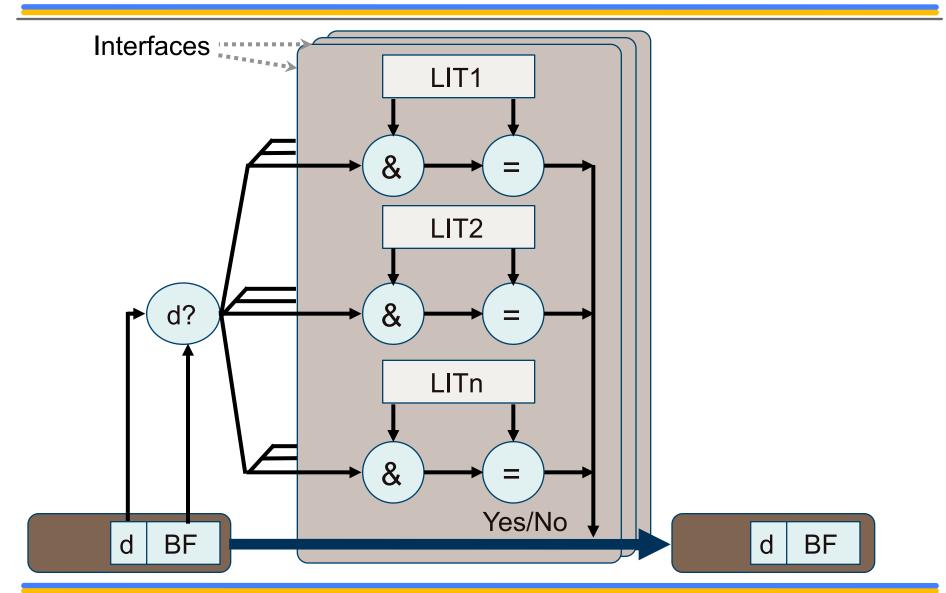
Using Link Identity Tags (LIT)

- Goal: Better false positive rate
 - Define n different LITs instead of a single LID
 - LIT has the same size as LID, and also k bits set to one
 - Power of choices
- Route creation and packet forwarding
 - Calculate n different candidate zFilters
 - Select the best performing zFilter (index d) and use that





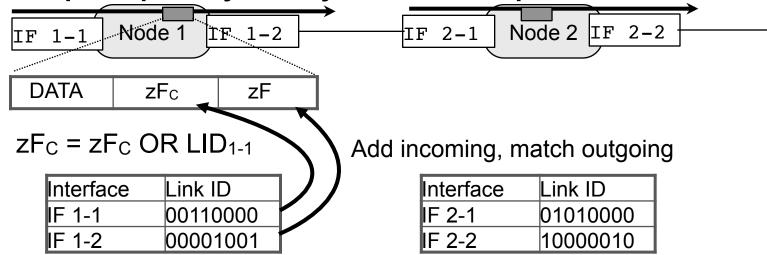
Using Link Identity Tags (LIT)





zFilter collection

- During packet traversal, the reverse zFilter can be easily generated
 - Add a field in the packet for collected zF
 - All routers forwarding the packet add the incoming LID to the field
 - Once the packet arrives to the destination, the collected zF can be used to forward data to the reverse direction
 - Simple especially with symmetric links/paths





Evaluation



Forwarding speed

- Measured on a NetFPGA
- Results
 - No routing table lookups
 → lower latency compared to IP
 - zF latency stays constant, independent of the network size
 - Line speed

Path	Avg. latency	Std dev.
Plain wire	94 µs	28 µs
IP router	102 μs	44 µs
zFilter	96 µs	28 µs

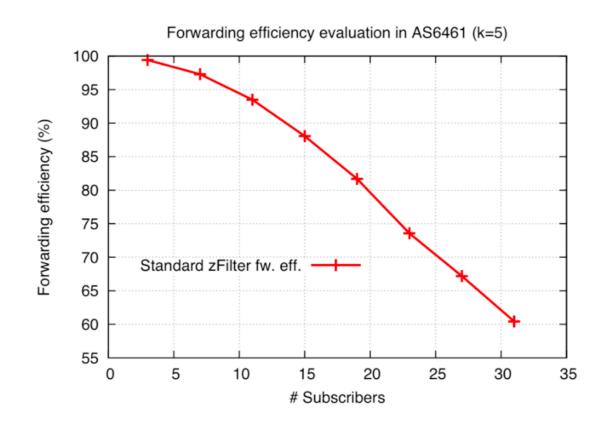
- Measurements in Blackadder (software)
 - Early results indicate that line speed forwarding over
 10 Gbit/s links can be achieved



Forwarding efficiency

- Simulations (ns-3) with
 - Rocketfuel
 - SNDlib
- Forwarding efficiency with 20 subscribers
 - **-~80%**
- AS6461:

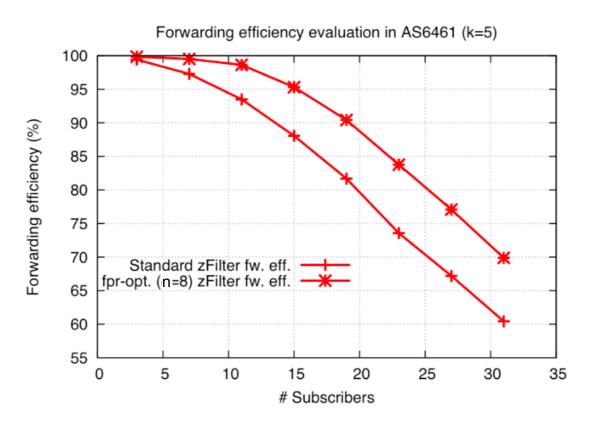
 138 nodes,
 372 links





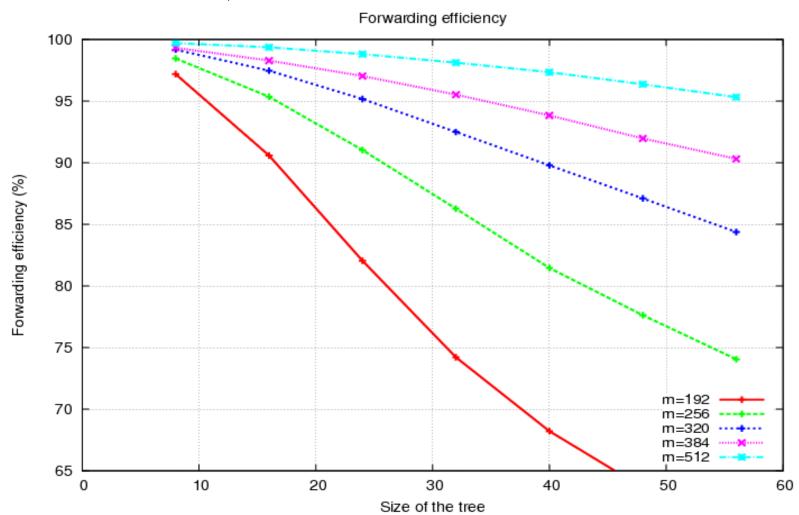
Forwarding efficiency

- Simulations with
 - Rocketfuel
 - SNDlib
- Forwarding efficiency with 20 subscribers
 - **-~80%**
 - LIT Optimized:88%



Changing zFilter size

AS3967: 79 nodes, 147 bi-directional links





Security

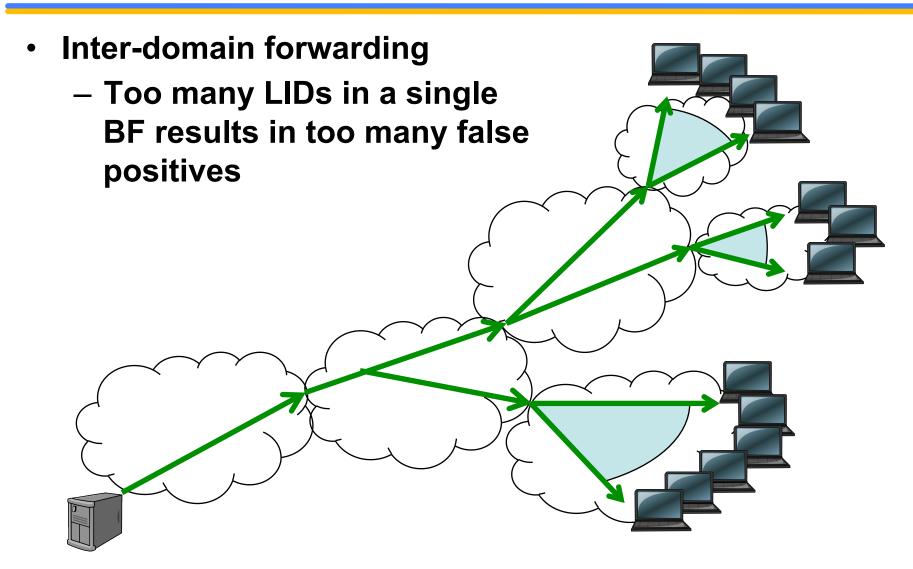
- A zFilter to a destination only works on a certain path, while IP addresses work from any source anywhere
 - → Better (although not complete) DDoS resistance
- zFilter doesn't reveal (directly) which nodes are involved in the communication
 - → Better privacy



Scalability enhancements

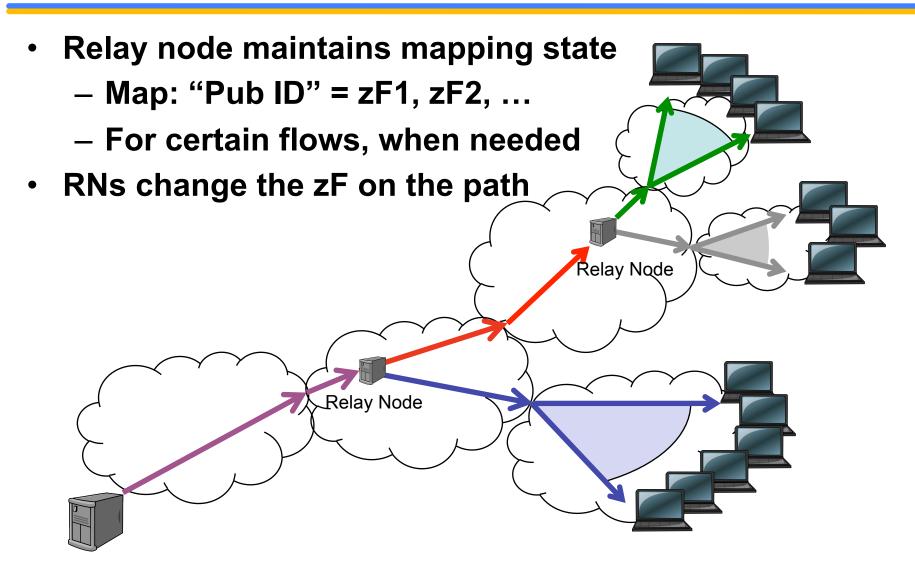


Scalability issues



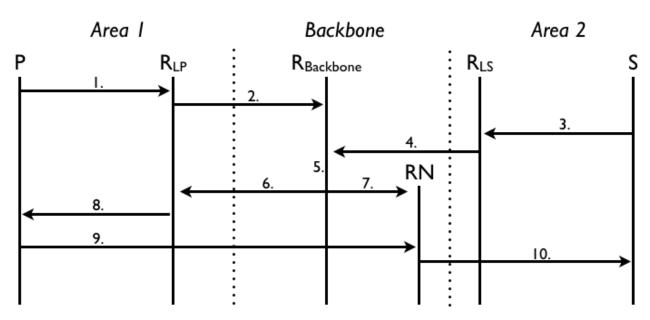


Scalability: Relay Nodes





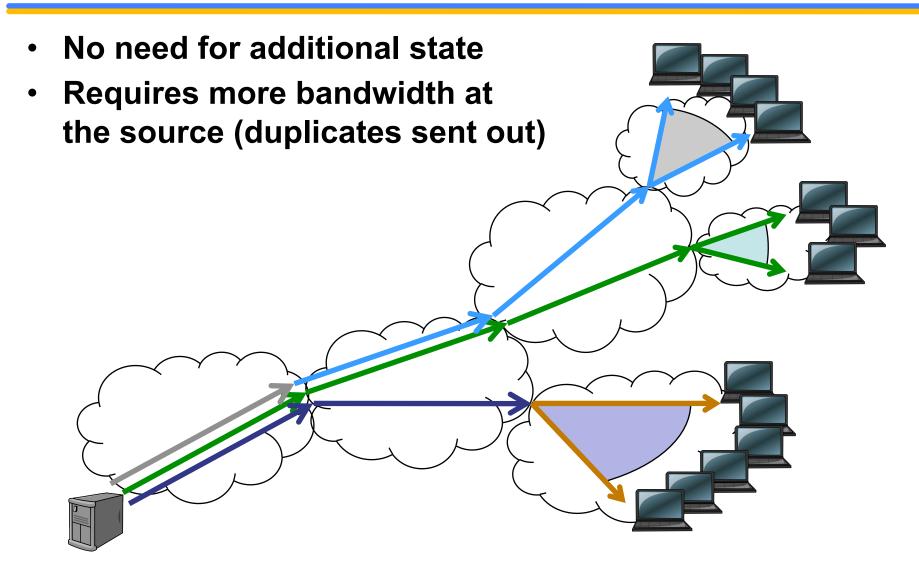
Setting up Relay Nodes



- I. & 2. Publish(Pub ID)
- 3. Subscribe(Pub ID)
- 4. Subscribe(Pub ID, BF_{area2}, NofLinks, NofSubscribers)
- 5. Decision, create a new Relay point
- 6. Subscribe(Pub ID, BF_{area1-to-RN}, NofLinks, NofSubscribers)
- 7. CreateRN(Pub ID, BF_{area2+RN-to-area2})
- 8. Subscribe(Pub ID, BF_{areal+areal-to-RN})
- 9. Deliver Data (Pub ID) to BF_{area I+area I-to-RN}
- 10. Deliver Data (Pub ID) to BF_{area2+RN-to-area2}



Scalability: Splitting the tree





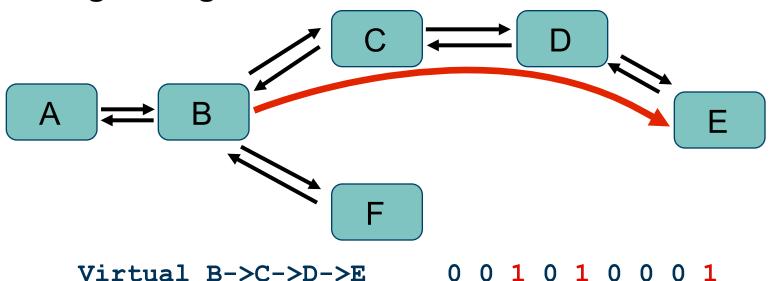
Scalability – stacking Bloom filters

- TM divides delivery tree into multiple parts along the paths
- Each part has its own BF
- These BFs are stacked into a packet, removed at boundaries
- BFs are variable size, chosen so that the probability of false positives is minimized



Scalability: Virtual trees

- Popular paths can be merged into virtual trees
 - A single Link ID for the tree
 - Additional state in the forwarding nodes
 - Increase scalability
- A virtual tree is not bound to a certain publication
 - E.g. a single tree for all AS transit traffic





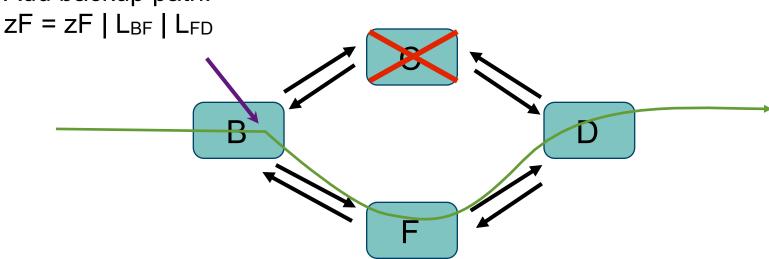
Failover enhancements



Fast reroute - Backup path

- Node B maintains backup path information
- In case of broken link, add backup path
 - Increases temporarily the false positive probability until a new path is calculated at the topology manager
 - No additional signaling

Add backup path:

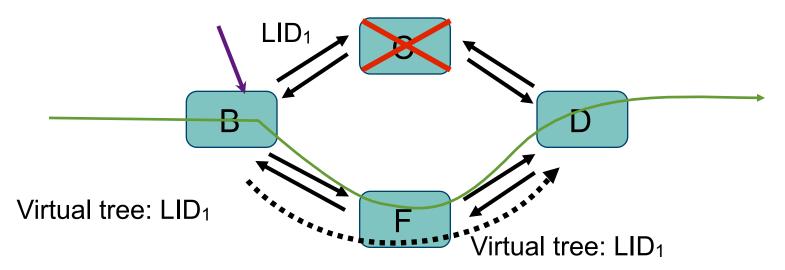




Fast reroute - Virtual trees

- zFilter unmodified
- Activate backup path in case of node failure
 - Adds signaling

Link broken, signal the activation of the backup path to F



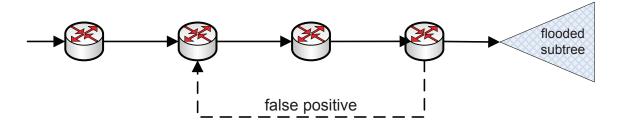


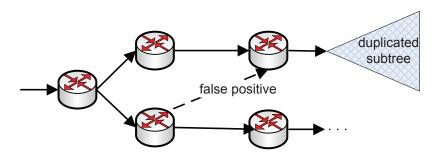
Loop prevention enhancements



Forwarding anomalies

- E.g. packet storms, forwarding loops, and flow duplication
- Accidental or maliciously created

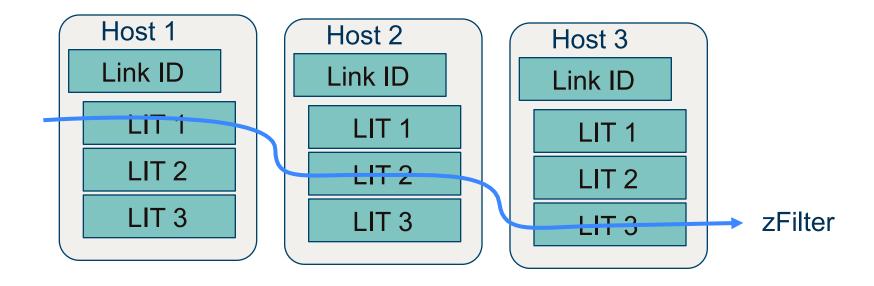






Avoiding loops

- Instead of fixed d determining the used LIT, change the d e.g. with d=(d+1) MOD e
- In case of a loop, the packet will have the same d only if the loop is e hops long
- Simple, stateless solution





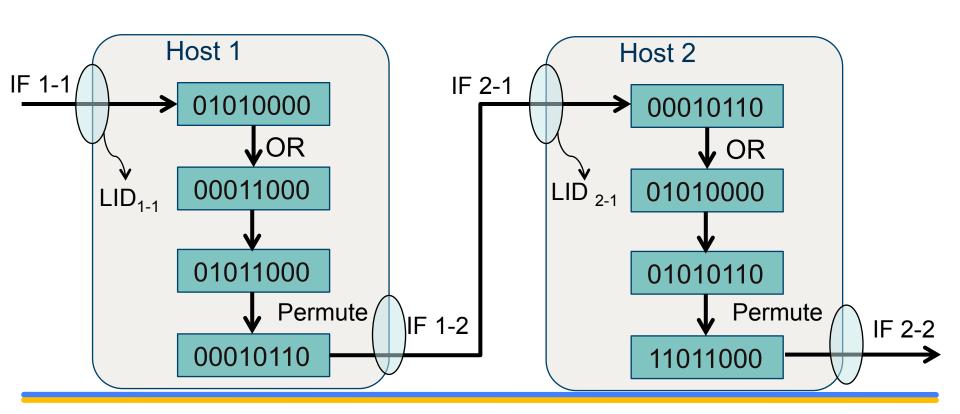
Permutations

- Goal: Prevent forwarding loops and flow duplication
- Idea: Make forwarding decision depend on the packet's path and hop-count
- Solution: Per-hop bit permutation
 - "Mix" the BF bits in incoming packets according to a function specific to the incoming interface
 - Simple to implement, no additional space in the packet, randomizes the BF in case of false positives
 - Requirements
 - Reversible operation, no significant increase in number of 1-bits
- Multicast zFilters
 - ORing is not enough, must be computed from the leaves of the tree



Forming the permuted zFilter

- Especially suitable when zF is collected on reverse path
- zFilter verification to the other direction
 - Permute with the function
 - Match outgoing interfaces





Security enhancements



Security weaknesses with static LID/LITs

- zFilter replay attacks
 - Sending data with the same zFilter
- Traffic injection attack
 - Using existing zFilter, send data from the middle of the path
- Computational attack
 - Collect zFilters from packets
 - Correlate zFilters to learn link IDs

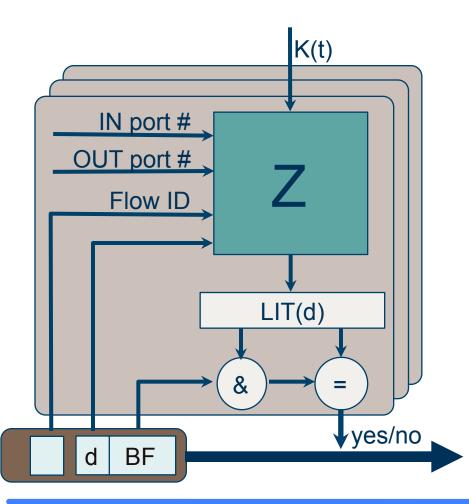


Secure forwarding

- Goal: to ensure (probabilistically) that hosts cannot send un-authorized traffic
- Solution (z-Formation): Compute LIT in line speed and bind it to
 - path: in-coming and out-going port
 - time: periodically changing keys
 - flow: flow identifier (e.g. content ID)



Secure case: z-Formation



- Form LITs algorithmically
 - at packet handling time
 - LIT(d) = Z(I,K(t),In,Out,d)
- Secure periodic key K
- Input port index
- Output port index
- Flow ID from the packet, e.g.
 - Information ID
 - IP addresses & ports
- d from the packet



Security properties

- Binding a zFilter only to the outgoing port
 - Traffic injection possible
 - Correlation attacks possible
- Bind to the incoming and outgoing ports
 - Traffic injection difficult (due to binding to incoming port)
 - Very hard to construct one without knowing keys along the path
 - Correlation attacks possible only for a given flow ID
 - Bound to the packet stream (flow ID)
- Need a cryptographically good Z-algorithm

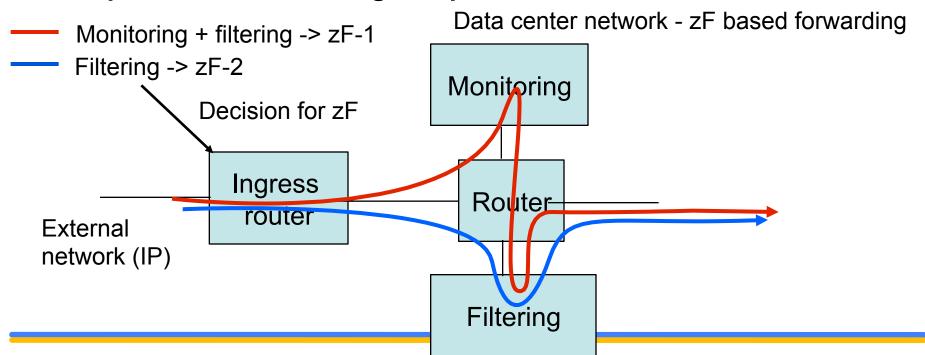


Applications



Data centers

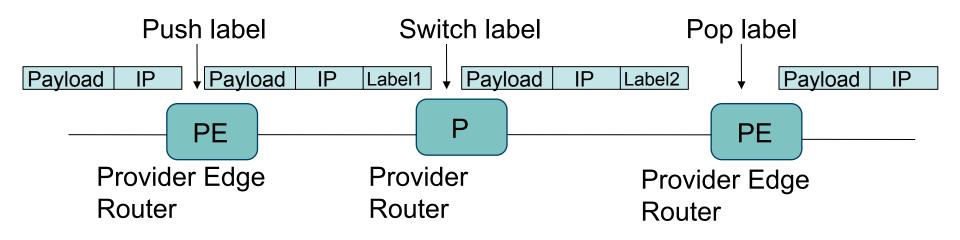
- zFilters only in the internal network
- Easier to modify the routing in the network
 - E.g. route packets via certain services: Load balancing, monitoring...
 - Binding the flow to input and output ports allows flexible path control at the ingress point





Background: (G)MPLS Multiprotocol label switching

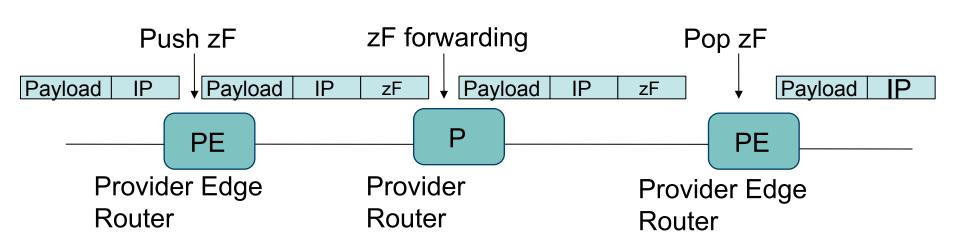
- Evolution: MPLS->MPLS-TE -> GMPLS
- (G)MPLS is a rich set of protocols
 - Setting up Label Switched Paths
 - Forwarding on the Label Switched Paths
 - Traffic Engineering, resiliency (e.g. fast reroute)
 - Enabler of VPN services
 - Control plane for many different technologies





MPSS Multiprotocol stateless switching

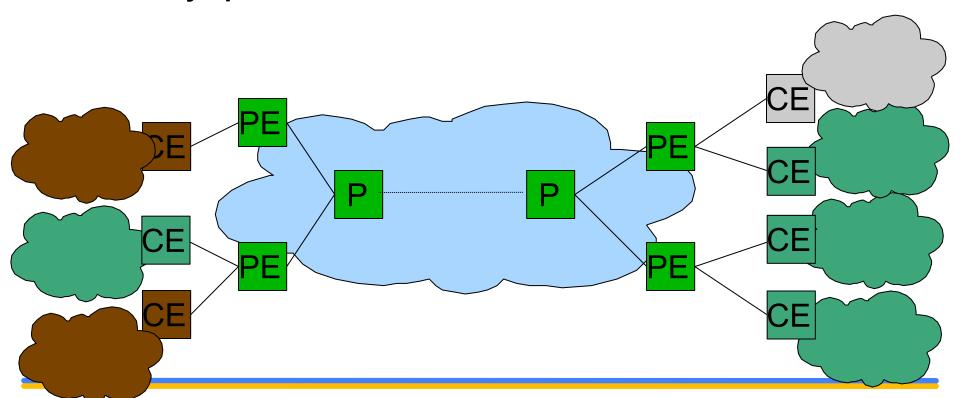
- Advantages over label switching
 - There is not necessarily need for signaling
 - In simpler case, no state required
 - Multicast support (setup, maintenance) much simpler than with (G)MPLS





Multicast VPN with MPSS

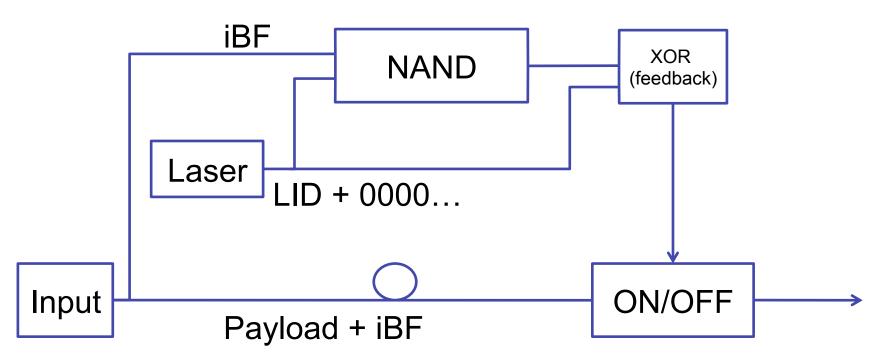
- Effective support of point-to-multipoint communication
- The bandwidth efficiency vs. Multicast state trade-off eliminated
 - (Though longer header sizes)
- With zFilters: no multicast states, and acceptable bandwidth efficiency up to ~20 PEs





Optical packet forwarding

- All-optical packet forwarding
 - Matching LID from laser with the iBF
 - XOR with feedback goes to zero if one bit fails (= no forwarding)





- New multicast forwarding mechanism
 - Suits pub/sub and synchronous multicast very well
 - Can also be applied outside our pub/sub model
 - Almost stateless
 - Good security properties
- But: Some scalability issues especially due to false positives
 - And also some security issues
- Many enhancements/changes/additions to the basic LIPSIN mechanism have been proposed
 - Tradeoffs
- E.g., work on inter-domain forwarding ongoing



Some references

- Petri Jokela, András Zahemszky, Christian Esteve, Somaya Arianfar, and Pekka Nikander, "*LIPSIN: Line speed Publish/Subscribe Inter-Networking*", ACM SIGCOMM 2009
- András Zahemszky and Somaya Arianfar, "Fast reroute for stateless multicast", IEEE RNDM 2009
- Christian Esteve et al., "Self-routing Denial-of-Service Resistant Capabilities using In-packet Bloom Filters", EC2ND, 2009
- Christian Esteve et al., "Data center networking with in-packet Bloom filters", SBRC, 2010
- András Zahemszky et al. "MPSS: Multiprotocol Stateless Switching", IEEE Global Internet Symposium 2010
- Mikko Särelä et al., "Forwarding Anomalies in Bloom Filter Based Multicast", IEEE INFOCOM 2010
- Dirk Trossen et al., "PURSUIT Deliverable D2.2: Conceptual Architecture: Principles, patterns and sub-components descriptions", Section 4.3: Forwarding, 2011
- Sajjad Rizvi, "Performance analysis of bloom filter-based multicast", Master's thesis, Aalto university, 2011