

Host Identity Protocol

Prof. Sasu Tarkoma

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Part of the material is based on lecture slides by Dr. Pekka Nikander (HIP) and Dmitrij Lagutin (PLA)

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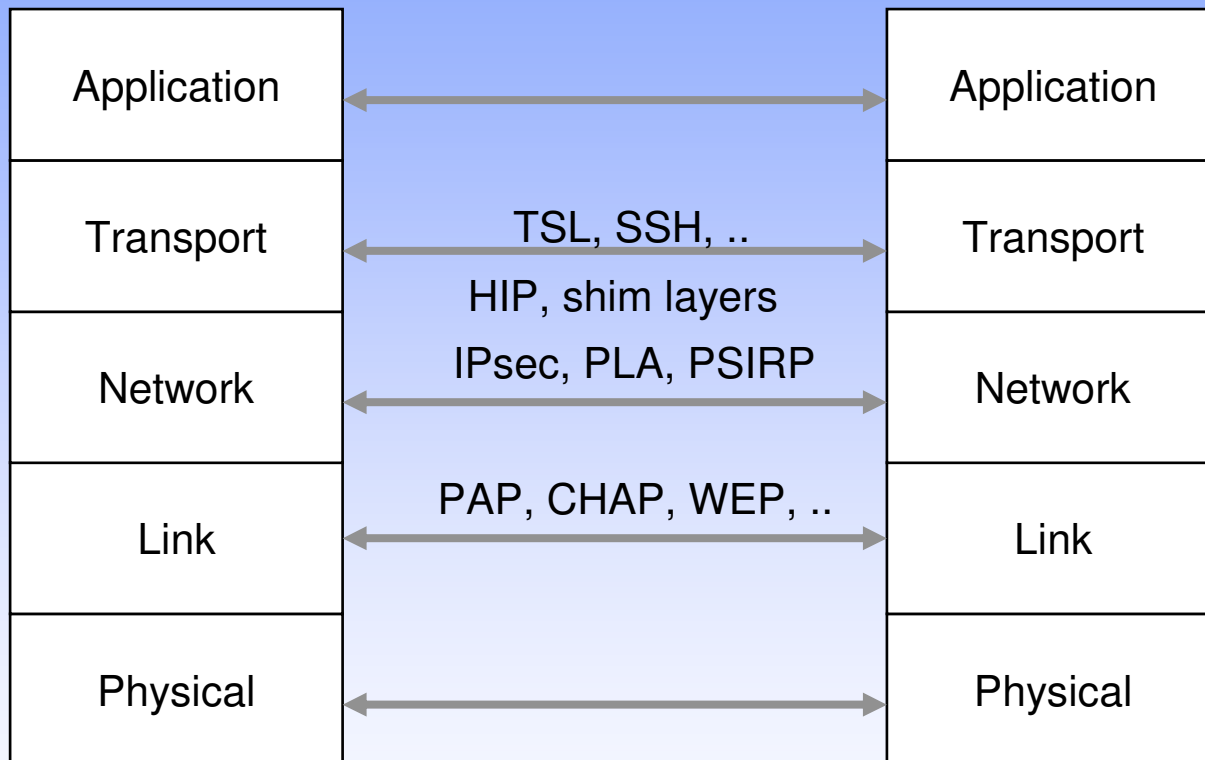
Introduction

- Current Internet is increasingly data and content centric
- The protocol stack may not offer best support for this
- End-to-end principle is no longer followed
 - Firewalls and NAT boxes
 - Peer-to-peer and intermediaries
- Ultimately, hosts are interested in receiving valid and relevant information and do not care about IP addresses or host names
- This motivate the design and development of new data and content centric networking architectures
 - Related work includes ROFL, DONA, TRIAD, FARA, AIP, ..

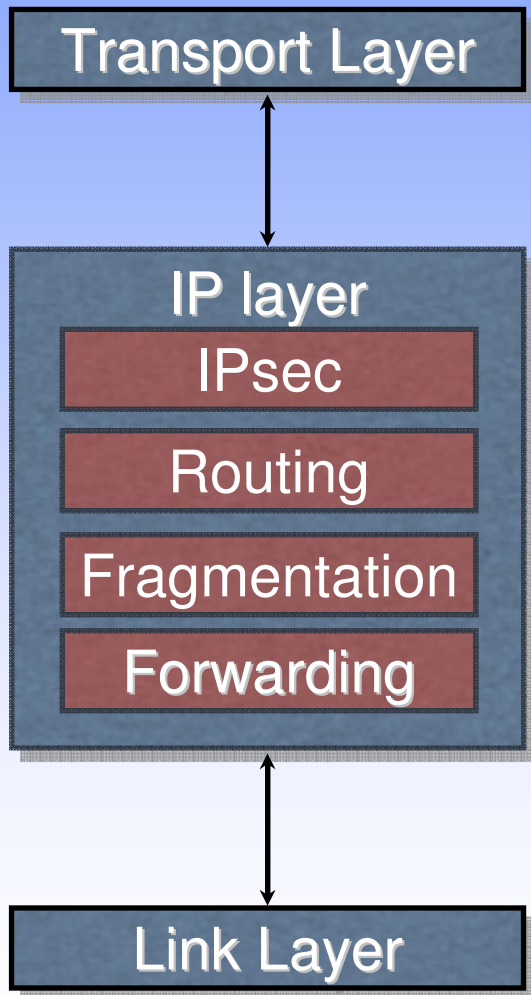
The Internet has Changed

- A lot of the assumptions of the early Internet has changed
 - Trusted end-points
 - Stationary, publicly addressable addresses
 - End-to-End
- We will have a look at these in the light of recent developments
- End-to-end broken by NATs and firewalls

HTTPS, S/MIME, PGP, WS-Security, Radius, Diameter, SAML 2.0 ..



Current State



Observations

End-to-end reachability is broken

Unwanted traffic is a problem

Mobility and multi-homing are challenging

Multicast is difficult (does not scale)

Security is difficult

Not optimal fit for broadcast and all-optical networking

HIP

What is HIP?

- HIP = Host Identity Protocol
- A proposal to separate identifier from locator at the network layer of the TCP/IP stack
 - A new name space of public keys
 - A protocol for discovering and authenticating bindings between public keys and IP addresses
- Secured using signatures and keyed hashes (hash in combination with a secret key)

Motivation

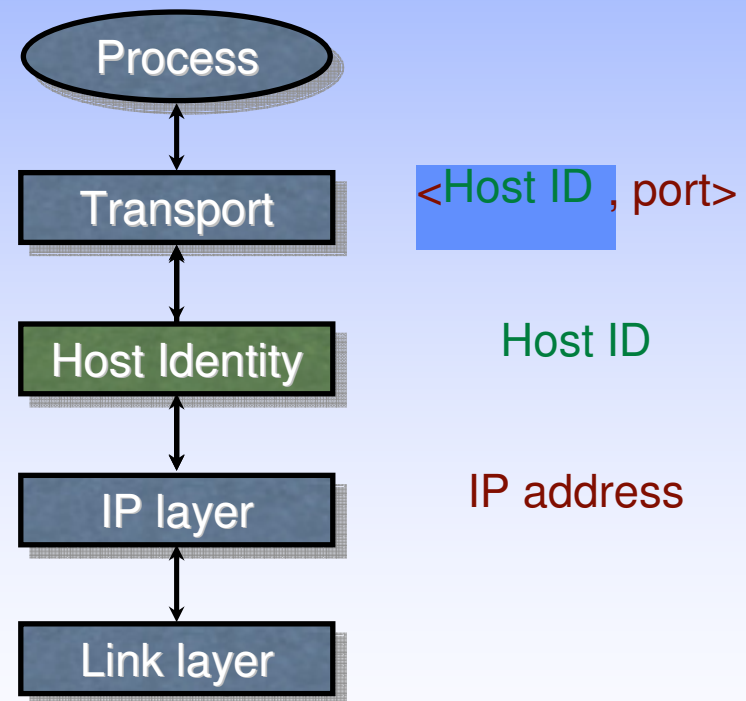
- Not to standardise a solution to a problem
 - No explicit problem statement
- Exploring the consequences of the id / loc split
 - Try it out in real life, in the live Internet
- A different look at many problems
 - Mobility, multi-homing, end-to-end security, signalling, control/data plane separation, rendezvous, NAT traversal, firewall security, ...

HIP in a Nutshell

- Architectural change to TCP/IP structure
- Integrates security, mobility, and multi-homing
 - Opportunistic host-to-host IPsec ESP
 - End-host mobility, across IPv4 and IPv6
 - End-host multi-address multi-homing, IPv4/v6
 - IPv4 / v6 interoperability for apps
- A new layer between IP and transport
 - Introduces cryptographic Host Identifiers

The Idea

- A new Name Space of Host Identifiers (HI)
 - Public crypto keys!
 - Presented as 128-bit long hash values, Host ID Tags (HIT)
- Sockets bound to HIs, not to IP addresses
- HIs translated to IP addresses in the kernel



Protocol overview

Initiator

Responder

I1: HIT_I , HIT_R or NULL

R1: HIT_I , [HIT_R , puzzle, DH_R , HI_R] _{sig}

I2: [HIT_I , HIT_R , solution, DH_I , { HI_I }] _{sig}

R2: [HIT_I , HIT_R , authenticator] _{sig}

User data messages

Control

Data

Base exchange

- Based on SIGMA family of key exchange protocols

standard authenticated Diffie-Hellman key exchange for session key generation

Initiator

I1 HIT_I, HIT_R or NULL

R1 HIT_I, [HIT_R

I2 [HIT_I, HIT_R

R2 [HIT_I, HIT_R

Select precomputed R1. Prevent DoS. Minimal state kept at responder! Does not protect from replay attacks.

verify, authenticate, replay protection

solve puzzle

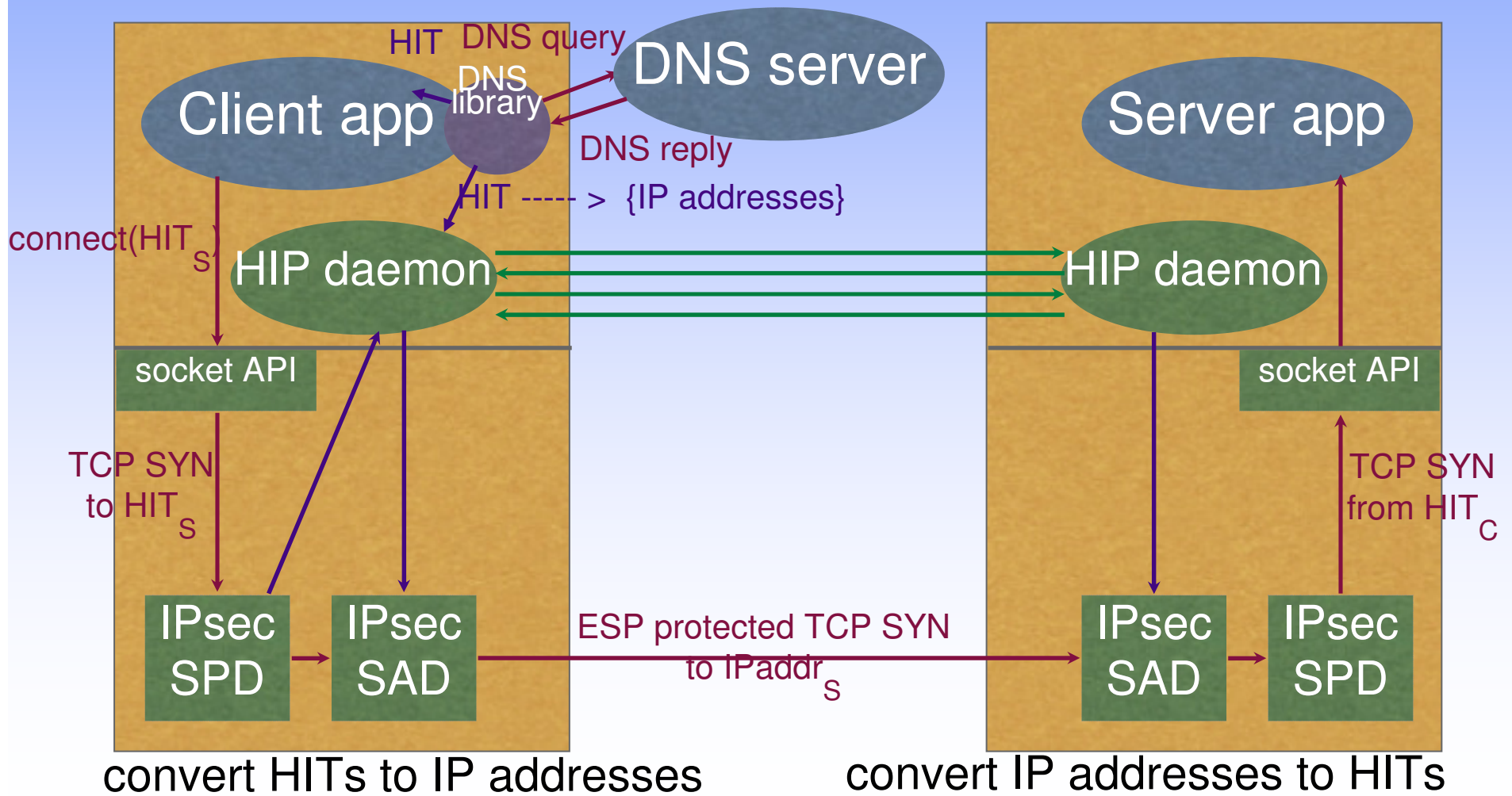
User data messages

ESP protected TCP/UDP, no explicit HIP header

Other Core Components

- Per-packet identity context
 - Indirectly, through SPI if ESP is used
 - Directly, e.g., through an explicit shim header
- A mechanism for resolving identities to addresses
 - DNS-based, if FQDNs used by applications
 - Or distributed hash tables (DHTs) based

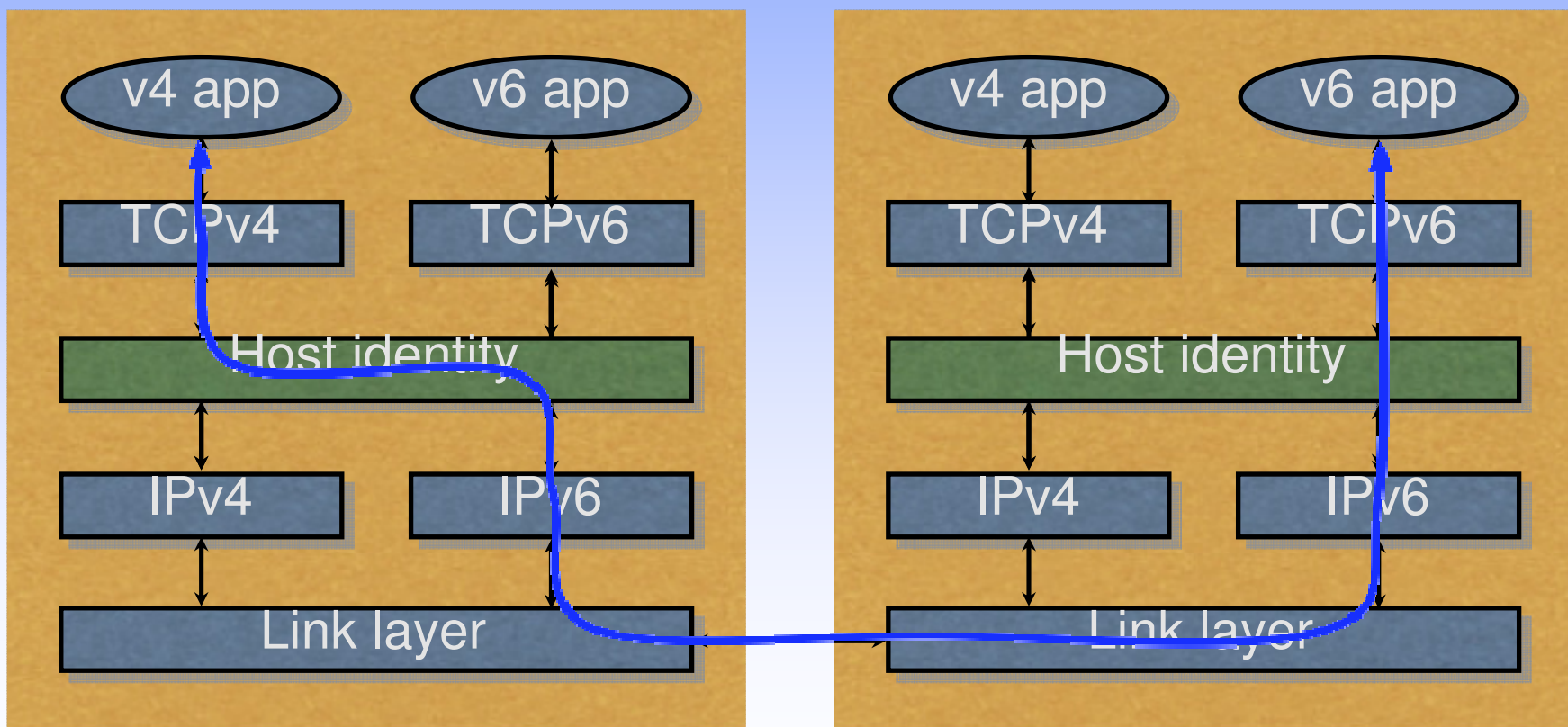
Using HIP with ESP



Many Faces

- More established views:
 - A different IKE for simplified end-to-end ESP
 - Super Mobile IP with v4/v6 interoperability and dynamic home agents
 - A host multi-homing solution
- Newer views:
 - New waist of IP stack; universal connectivity
 - Secure carrier for signalling protocols

HIP as the new waist of TCP/IP

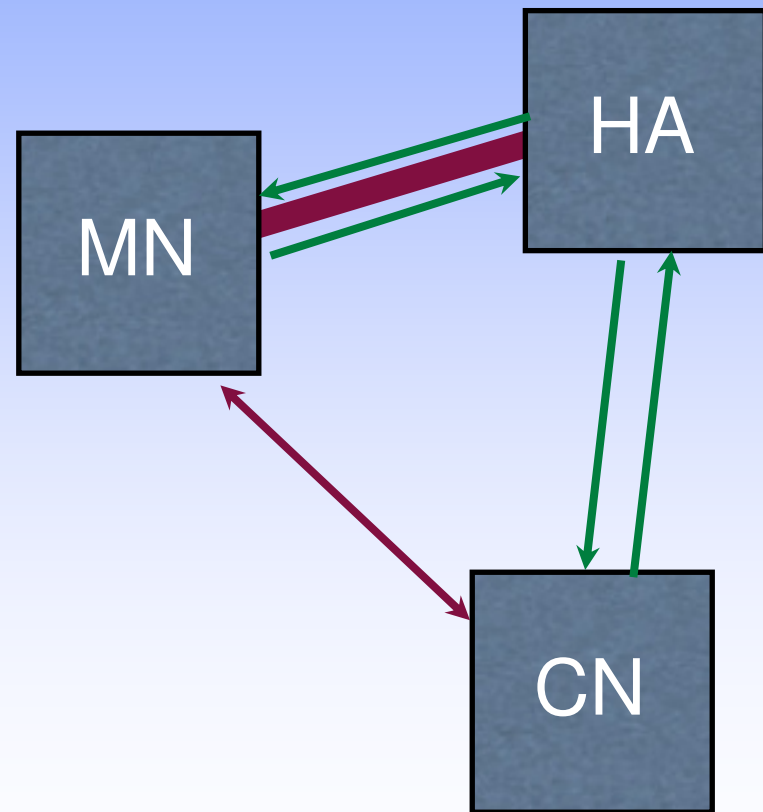


Rendezvous

- Initial rendezvous
 - How to find a moving end-point?
 - Can be based on directories
 - Requires fast directory updates
 - Bad match for DNS
- Tackling double-jump
 - What if both hosts move at same time?
 - Requires rendezvous point

Mobile IP

- Home Agent (HA)
 - Serves a Home Address
 - Initial reachability
 - Triangular routing
- Route optimization
 - Tunnels to bypass HA
 - HA as rendezvous point



Mobility protocol

Mobile

Corresponding

UPDATE: HITs, new locator(s), sig

UPDATE: HITs, RR challenge, sig

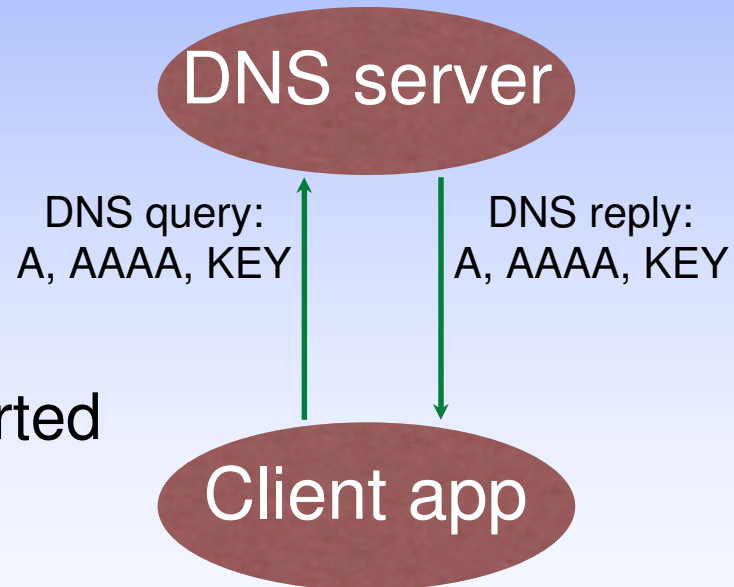
ESP from MN to CN

UPDATE: HITs, RR response, sig

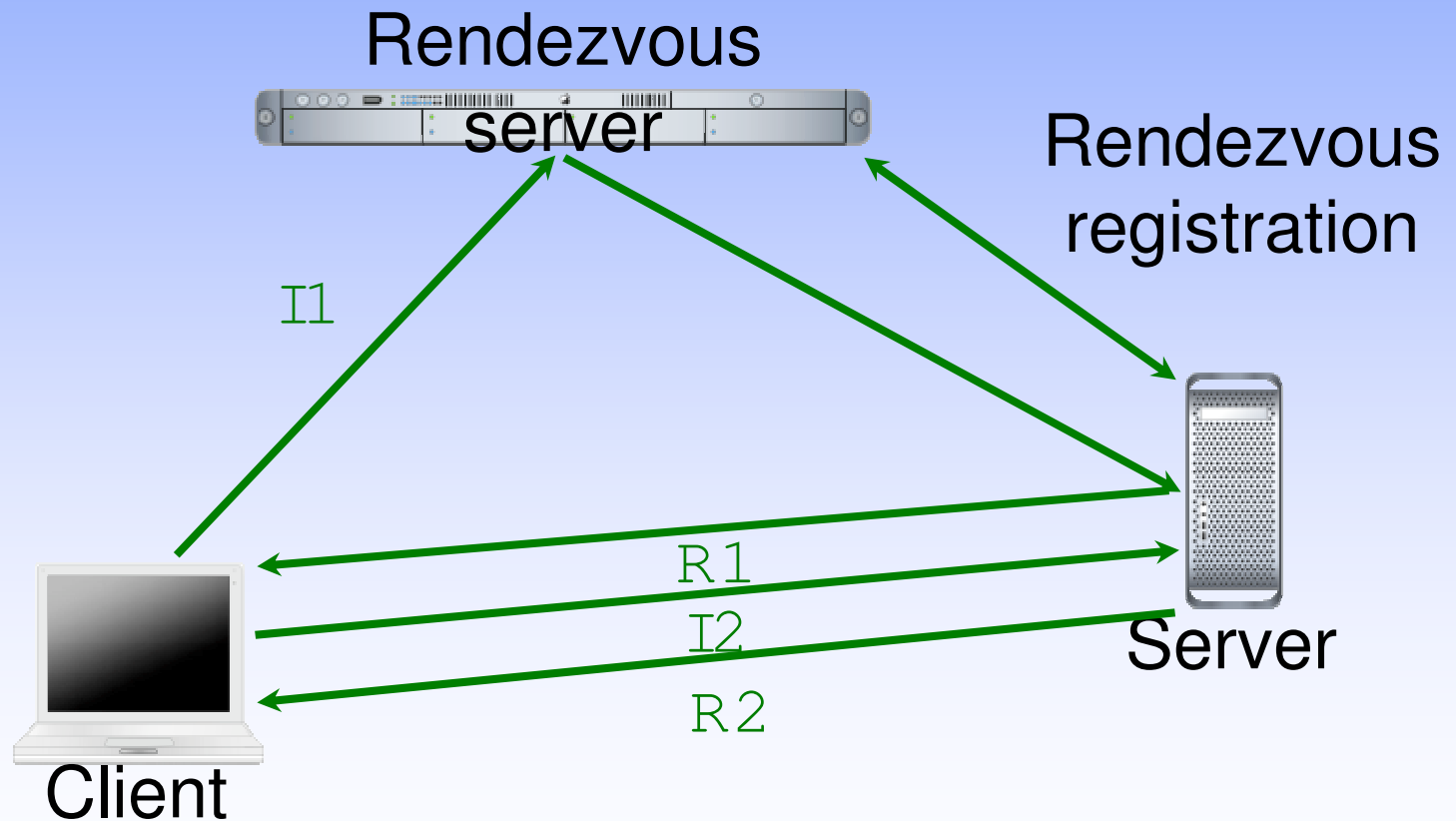
ESP on both directions

Key distribution for HIP

- Depends on application
- For multi-addressing, self-generated keys
- Usually keys in the DNS
- Can use PKI if needed
- Opportunistic mode supported
 - SSH-like leap-of-faith
 - Accept a new key if it matches a fingerprint



Basic HIP rendezvous



The infrastructure question

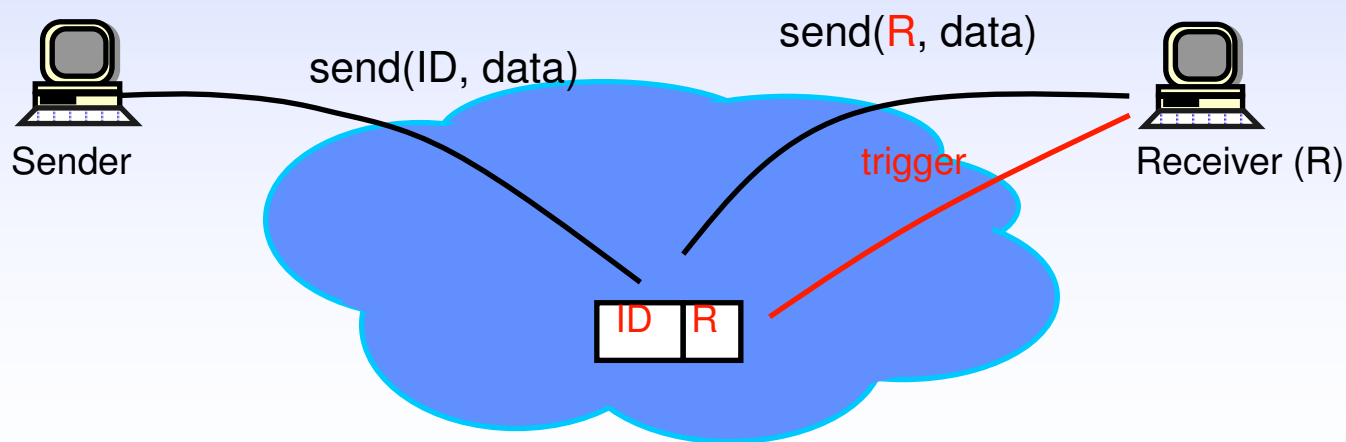
- HIs originally planned to be stored in the DNS
 - Retrieved simultaneously with IP addresses
 - Does not work if you have only a HIT
- Question: How to get data based on HIT only?
 - HITs look like 128-bit **random** numbers
- Possible answer: DHT based overlay like i^3

Distributed Hash Tables

- Distributed directory for flat data
- Several different ways to implement
- Each server maintains a partial map
- Overlay addresses to direct to the right server
- Resilience through parallel, unrelated mappings
- Used to create **overlay networks**

i^3 rendezvous abstraction

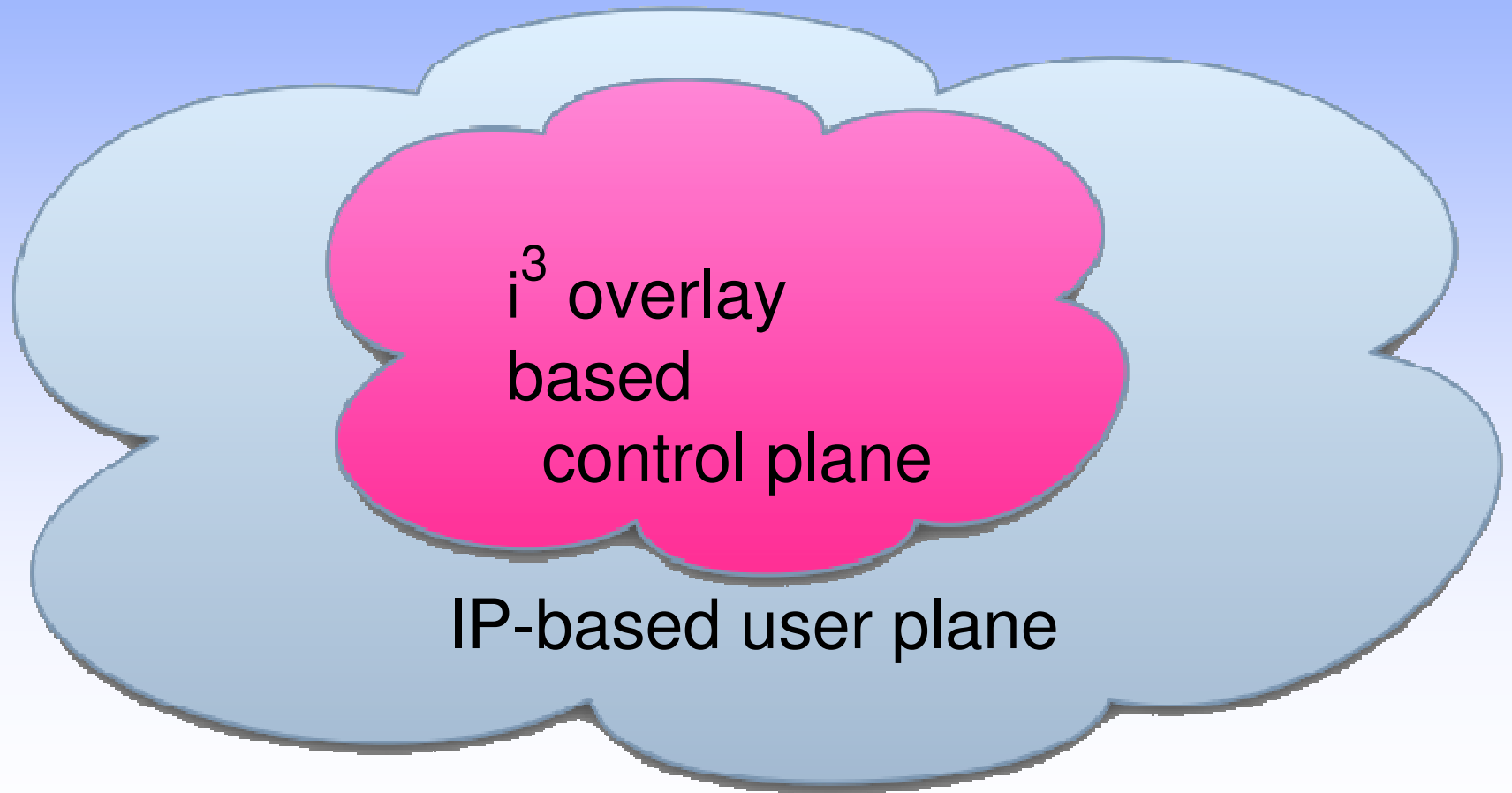
- Trigger inserted by receiver(s)
- Packets addressed to identifiers
- i^3 routes packet to the receiver(s)



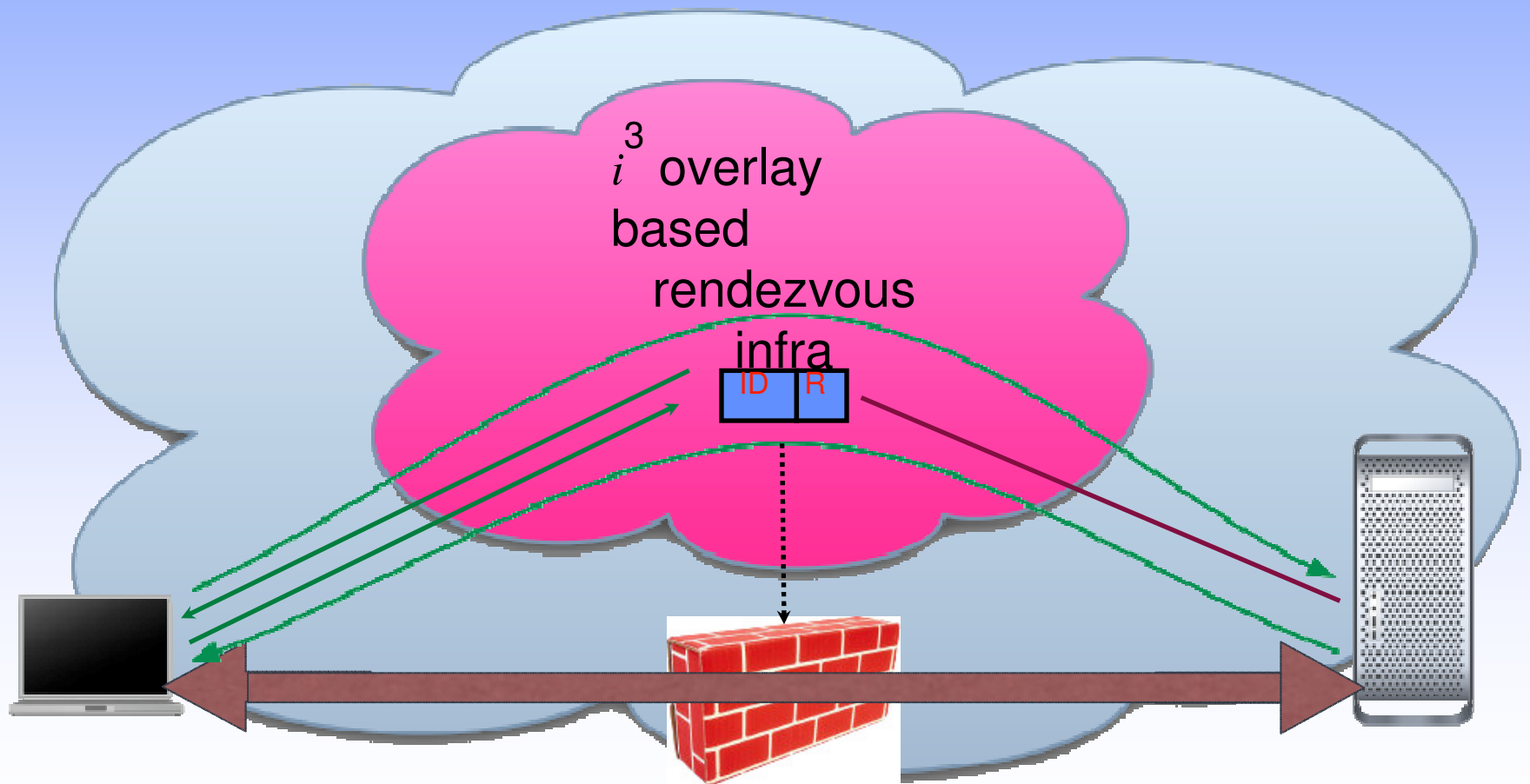
Hi³: combining HIP and i3

- Developed at Ericsson Research IP Networks
- Uses i³ overlay for HIP *control* packets
 - Provides rendezvous for HIP
- *Data* packets use plain old IP
 - Cryptographically protected with ESP
- Only soft or optional state in the network

H_i^3 and DHT-based rendezvous



Control/data separation



An Internet control plane?

- HIP separates control and data traffic
- Hi³ routes control traffic through overlay
 - Control and data packets take potentially very different paths
- *Allows* telecom-like control ...
 - ... but does not *require* it

Current status

- RFCs

- Host Identity Protocol (HIP) Architecture (RFC 4423) (60977 bytes)
- Host Identity Protocol (RFC 5201) (240492 bytes)
- Host Identity Protocol (HIP) Domain Name System (DNS) Extensions (RFC 5205) (34799 bytes)
- Host Identity Protocol (HIP) Registration Extension (RFC 5203) (26620 bytes)
- Using the Encapsulating Security Payload (ESP) Transport Format with the Host Identity Protocol (HIP) (RFC 5202) (68195 bytes)
- Host Identity Protocol (HIP) Rendezvous Extension (RFC 5204) (30233 bytes)
- End-Host Mobility and Multihoming with the Host Identity Protocol (RFC 5206) (99430 bytes)
- Using the Host Identity Protocol with Legacy Applications (RFC 5338) (34882 bytes)

- Internet-Drafts

- Basic HIP Extensions for Traversal of Network Address Translators (75933 bytes)
- Basic Socket Interface Extensions for Host Identity Protocol (HIP) (42500 bytes)
- HIP Certificates (19638 bytes)
- HIP BONE: Host Identity Protocol (HIP) Based Overlay Networking Environment (42692 bytes)

Implementations

- HIP for Linux (HIPL) infrachip.hiit.fi
- HIP for NetBSD
- OpenHIP